

The RDF language

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Center for Web Research
<http://www.cwr.cl>

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RDF Language and Model Theory looks a little bit complicated at first sight...



RDF Semantics

W3C Recommendation 10 February 2004

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Please refer to the [errata](#) for this document, which may include some norm

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RDF Semantics

W3C Recommendation 10

This Version:

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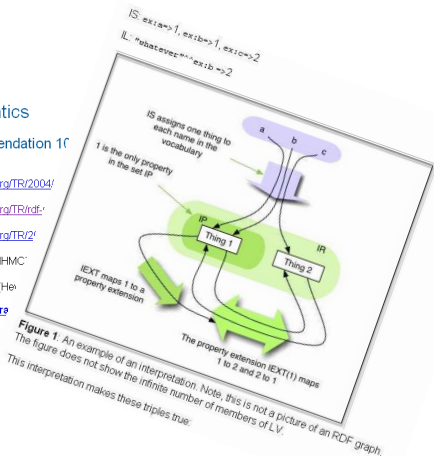
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RDF Language and Model Theory looks a little bit complicated at first sight...

$\text{IEXT}(\{(rdfs:subPropertyOf)\})$ is transitive and reflexive on IP

If $\langle x, y \rangle$ is in $\text{IEXT}(\{(rdfs:subPropertyOf)\})$ then x and y are in IP and $\text{IEXT}(x)$ is a subset of $\text{IEXT}(y)$

If x is in IC then $\langle x, \{(rdfs:Resource)\} \rangle$ is in $\text{IEXT}(\{(rdfs:subClassOf)\})$

If $\langle x, y \rangle$ is in $\text{IEXT}(\{(rdfs:subClassOf)\})$ then x and y are in IC and $\text{ICEXT}(x)$ is a subset of $\text{ICEXT}(y)$

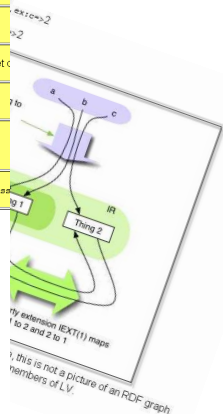
$\text{IEXT}(\{(rdfs:subClassOf)\})$ is transitive and reflexive on IC

If x is in $\text{ICEXT}(\{(rdfs:ContainerMembershipProperty)\})$ then:
 $\langle x, \{(rdfs:member)\} \rangle$ is in $\text{IEXT}(\{(rdfs:subPropertyOf)\})$

If x is in $\text{ICEXT}(\{(rdfs:Datatype)\})$ then $\langle x, \{(rdfs:Literal)\} \rangle$ is in $\text{IEXT}(\{(rdfs:subClassOf)\})$

RDFS axiomatic triples.

```
rdfs:type rdfs:domain rdfs:Resource .
rdfs:domain rdfs:domain rdfs:Property .
rdfs:range rdfs:domain rdfs:Property .
rdfs:subPropertyOf rdfs:domain rdfs:Property .
rdfs:subClassOf rdfs:domain rdfs:Class .
rdfs:subject rdfs:domain rdfs:Statement .
rdfs:predicate rdfs:domain rdfs:Statement .
rdfs:object rdfs:domain rdfs:Statement .
rdfs:member rdfs:domain rdfs:Resource .
rdfs:first rdfs:domain rdfs:List .
rdfs:rest rdfs:domain rdfs:List .
rdfs:seeAlso rdfs:domain rdfs:Resource .
rdfs:isDefinedBy rdfs:domain rdfs:Resource .
rdfs:comment rdfs:domain rdfs:Resource .
```



RDF Language and Model Theory looks a little bit complicated at first sight...

`!EXT((rdf:subPropertyOf))` is transitive and reflexive on IP
 If $\langle x, y \rangle$ is in `!EXT((rdf:subPropertyOf))` then x and y are in IP and `!EXT((rdf:subPropertyOf))` is a subset of IP
 If x is in IC then $\langle x, !\text{EXT}(\text{rdf:Resource}) \rangle$ is in `!EXT((rdf:subPropertyOf))`
 If $\langle x, y \rangle$ is in `!EXT((rdf:subPropertyOf))` then x and y are in IP and `!EXT((rdf:subPropertyOf))` is a subset of IP
 If x is in IC then $\langle x, !\text{EXT}(\text{rdf:Resource}) \rangle$ is in `!EXT((rdf:subPropertyOf))`

RDF entailment rules	
Rule Name	then add
rdf1	<code>aaa rdf:type rdf:Property</code> <code>!rnn rdf:type rdf:XMLLiteral</code> where <code>!rnn</code> identifies a blank node allocated to <code>!ll</code> by rule <code>lg</code>
rdf2	<code>!ll</code> is a well-typed XML literal where <code>!ll</code> is a well-typed XML literal

These rules are complete in the following sense.
RDF entailment lemma S *rdf-entails* E if and only if there is a graph which can be derived from S plus the RDF axiomatic triples by the application of rule `lg` and the RDF entailment rules and which simply entails E . (Proof in Appendix A)

Note that this does not require the use of rule `gl`.

7.3 RDFS Entailment Rules

RDFS entailment rules.	
Rule Name	then add:
rdfs1	<code>!ll</code> is a plain literal (with or without a language tag) where <code>!ll</code> is a plain literal (with or without a language tag)

`!rnn` *rdfs:type* `rdfs:Literal`
 where `!rnn` identifies a blank node allocated to `!ll` by rule `lg`

a picture of an RDF graph.

RDF Language and Model Theory looks a little bit complicated at first sight...

$\text{IEXT}(\text{rdf:type})$ is transitive and reflexive on IP
 If $\langle x, y \rangle$ is in $\text{IEXT}(\text{rdf:type})$ then x and y are in IP and $\text{IEXT}(\text{rdf:type})$ is a subset of $\text{IEXT}(\text{rdf:type})$
 If x is in IP
 If $\langle x, y \rangle$ is in $\text{IEXT}(\text{rdf:type})$ then y is in IP
 If $\langle x, y \rangle$ is in $\text{IEXT}(\text{rdf:type})$ then y is in IP

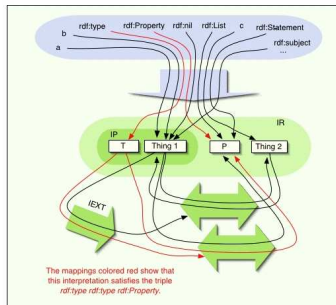


Figure 2: An rdf-interpretation.

This is not the smallest rdf-interpretation which extends the earlier example, since one made $\text{IEXT}(T)$ be $\{<1,2>, <T,2>\}$, and managed without having P in the universe. In general, the picture of an RDF graph.

Note
 7:

Rule Name	
rule1	<p> where ll is a plain literal (with or without a language tag). </p>

Several dozens of reserved keywords...

<code>rdfs:Resource [res]</code>	<code>rdf:type [type]</code>	<code>rdfs:isDefinedBy [isDefined]</code>
<code>rdf:Property [prop]</code>	<code>rdfs:domain [dom]</code>	<code>rdfs:comment [comment]</code>
<code>rdfs:Class [class]</code>	<code>rdfs:range [range]</code>	<code>rdfs:label [label]</code>
<code>rdfs:Literal [literal]</code>	<code>rdfs:subClassOf [sc]</code>	<code>rdf:value [value]</code>
<code>rdfs:Datatype [datatype]</code>	<code>rdfs:subPropertyOf [sp]</code>	<code>rdf:nil [nil]</code>
<code>rdf:XMLLiteral [xmlLit]</code>	<code>rdf:subject [subj]</code>	<code>rdf:_1 [_1]</code>
<code>rdfs:Container [cont]</code>	<code>rdf:predicate [pred]</code>	<code>rdf:_2 [_2]</code>
<code>rdf:Statement [stat]</code>	<code>rdf:object [obj]</code>	<code>...</code>
<code>rdf:List [list]</code>	<code>rdfs:member [member]</code>	<code>rdf:_i [_i]</code>
<code>rdf:Alt [alt]</code>	<code>rdf:first [first]</code>	<code>...</code>
<code>rdf:Bag [bag]</code>	<code>rdf:rest [rest]</code>	
<code>rdf:Seq [seq]</code>	<code>rdfs:seeAlso [seeAlso]</code>	
<code>rdfs:ContainerMembershipProperty [contMP]</code>		

...plus axiomatic triples...

(type,	type,	prop)
(subj,	type,	prop)
(pred,	type,	prop)
(obj,	type,	prop)
(first,	type,	prop)
(rest,	type,	prop)
(value,	type,	prop)
(_1,	type,	prop)
(_1,	type,	contMP)
(_2,	type,	prop)
(_2,	type,	contMP)
...		
(_i,	type,	prop)
(_i,	type,	contMP)
...		
(nil,	type,	prop)
(xmlLit,	type,	datatype)

...plus more axiomatic triples...

(type,	dom,	res)	(type,	range,	class)
(dom,	dom,	prop)	(dom,	range,	class)
(range,	dom,	prop)	(range,	range,	class)
(sp,	dom,	prop)	(sp,	range,	prop)
(sc,	dom,	class)	(sc,	range,	class)
(subj,	dom,	stat)	(subj,	range,	res)
(pred,	dom,	stat)	(pred,	range,	res)
(obj,	dom,	stat)	(obj,	range,	res)
(member,	dom,	res)	(member,	range,	res)
(first,	dom,	list)	(first,	range,	res)
(rest,	dom,	list)	(rest,	range,	list)
(seeAlso,	dom,	res)	(seeAlso,	range,	res)
(isDefined,	dom,	res)	(isDefined,	range,	res)
(comment,	dom,	res)	(comment,	range,	literal)
(label,	dom,	res)	(label,	range,	literal)
(value,	dom,	res)	(value,	range,	res)
(_1,	dom,	res)	(_1,	range,	res)
(_2,	dom,	res)	(_2,	range,	res)
...			...		
(_i,	dom,	res)	(_i,	range,	res)
...			...		

...plus more axiomatic triples...

```
(alt,      sc,      cont)
(bag,      sc,      cont)
(seq,      sc,      cont)
(contMP,   sc,      prop)
(xmlLit,   sc,      literal)
(datatype, sc,      class)

(isDefined, sp,      seeAlso)
```

...and on top of this a (slightly) non-standard model theory

- ▶ A notion of **interpretation**

$(Res, Prop, Class, Lit, PExt, CExt, Int)$

including subsets of the universe denoting properties, classes and literals, and mapping defining their extensions.

- ▶ Notion of interpretation of blank nodes
- ▶ Definition of reflexivity, transitivity and semi-extensionality of `subClass` and `subProperty`
- ▶ Typing restrictions

But...we need a workable language to bring to reality the vision of the Semantic Web

Would like to:

- ▶ Have a **simple user-language** to be able to popularize RDF among Web users.
- ▶ Have a **simple specification** to allow sound development work.
- ▶ Have a **language in streamlined form** to make it easy to formalize and prove results about its properties.

What is to be done?: To simplify the language

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There is a minimal fragment of the theory
preserving the essential core of RDFS

What is to be done?: To simplify the language



There is a minimal fragment of the theory preserving the essential core of RDFS

- ▶ Basic idea: Separate user-language from features and constructors which define and specify the language.
- ▶ Concentrate in vocabulary with non-trivial semantics.

Main contributions & Outline

- ▶ Identify a fragment of RDFS that covers the crucial vocabulary and preserves the original RDFS semantics.
- ▶ Study dependencies among vocabulary and develop sound and complete proof systems for each fragment.
- ▶ Present algorithms to modularize reasoning according to relevant vocabulary.

Main contributions & Outline

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- ▶ Present algorithms to modularize reasoning according to relevant vocabulary.

The core vocabulary

<code>rdfs:Resource [res]</code>	<code>rdf:type [type]</code>	<code>rdfs:isDefinedBy [isDefined]</code>
<code>rdf:Property [prop]</code>	<code>rdfs:domain [dom]</code>	<code>rdfs:comment [comment]</code>
<code>rdfs:Class [class]</code>	<code>rdfs:range [range]</code>	<code>rdfs:label [label]</code>
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$$pdf = \{sp, sc, type, dom, range\}$$

Blank node rule

$\frac{G}{H}$ if there is a homomorphism $\mu : H \rightarrow G$

Core rules

Subproperty (transitivity, definition)

$$\frac{(A, \text{sp}, B) \ (B, \text{sp}, C)}{(A, \text{sp}, C)}$$

$$\frac{(A, \text{sp}, B) \ (X, A, Y)}{(X, B, Y)}$$

Subclass (transitivity, definition)

$$\frac{(A, \text{sc}, B) \ (B, \text{sc}, C)}{(A, \text{sc}, C)}$$

$$\frac{(A, \text{sc}, B) \ (X, \text{type}, A)}{(X, \text{type}, B)}$$

Typing (domain, range)

$$\frac{(A, \text{dom}, B) \ (X, A, Y)}{(X, \text{type}, B)}$$

$$\frac{(A, \text{range}, B) \ (X, A, Y)}{(Y, \text{type}, B)}$$

Implicit Typing (strange case...)

$$\frac{(A, \text{dom}, B) \ (C, \text{sp}, A) \ (X, C, Y)}{(X, \text{type}, B)}$$

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$$\frac{(A, \text{range}, B) \ (C, \text{sp}, A) \ (X, C, Y)}{(Y, \text{type}, B)}$$

Reflexivity rules

If “ $(A, \text{type}, \text{Property})$ ” then (A, sp, A)

Subproperty Reflexivity

$$\frac{(X, A, Y)}{(A, \text{sp}, A)}$$

$$\frac{(A, \text{sp}, B)}{(A, \text{sp}, A) \ (B, \text{sp}, B)}$$

$$\overline{(p, \text{sp}, p)} \text{ for } p \in \rho_{df}$$

$$\frac{(A, \text{dom}, X)}{(A, \text{sp}, A)} \quad \frac{(A, \text{range}, X)}{(A, \text{sp}, A)}$$

Subclass Reflexivity

$$\frac{(A, \text{sc}, B)}{(A, \text{sc}, A) \ (B, \text{sc}, B)}$$

$$\frac{(X, \text{dom}, A)}{(A, \text{sc}, A)}$$

$$\frac{(X, \text{range}, A)}{(A, \text{sc}, A)}$$

$$\frac{(X, \text{type}, A)}{(A, \text{sc}, A)}$$

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Soundness and Completeness

Let \models denote the standard RDFS entailment, and $\vdash_{\rho\text{df}}$ a proof system based on the rules presented.

Theorem

Let G and H be graphs in ρdf then

$$G \models H \text{ if and only if } G \vdash_{\rho\text{df}} H.$$

Blank Nodes Modularization

Blank nodes can be treated in an orthogonal form to ρ df vocabulary.

Theorem

Let G and H be graphs in ρ df and $G \models H$, then

*there is a proof of H from G where the blank rule is used **at most once** and **at the end**.*

The role of reflexivity

The **only** consequence of reflexivity of `sp` and `sc` in RDFS is the possible entailment of triples of the form (x, sp, x) , (x, sc, x) .

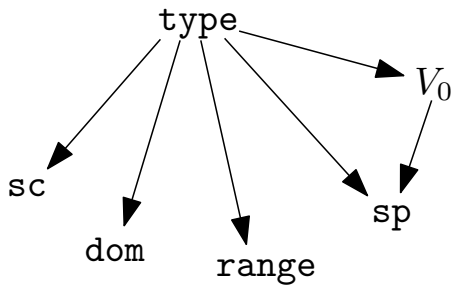
Theorem

Let G and H be ρ df graphs. Assume that H does not contain triples of the form (x, sp, x) and (x, sc, x) . Then,

$G \vdash_{\rho\text{df}} H$ without using reflexivity rules.

(Also, by not imposing reflexivity, axiomatic triples can be completely avoided.)

Dependence diagram among ρ df vocabulary



To determine $G \models H$ it is enough to test $G' \models H$ where G' is the subgraph of G which involves **only** nodes in $\text{voc}(H)$ and their dependencies in the diagram.

It is possible to avoid the closure to test RDFS entailment

- ▶ A naive approach to test $G \models H$ is:
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The size of the closure of G is $O(|G|^2)$, and this bound is tight.

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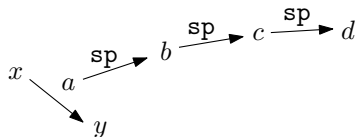
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- ▶ Alternative: to use a **goal oriented** approach based on the dependencies diagram.

Goal oriented entailment algorithm

Does the graph entails (x, d, y) ?

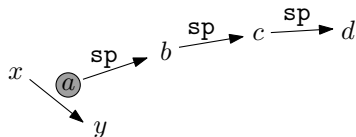
Look for triples of the form (x, a, y) and sp-paths from a to d .



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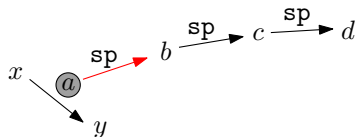
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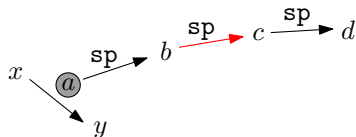
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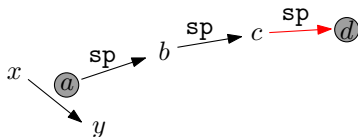
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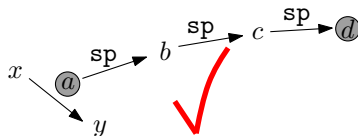
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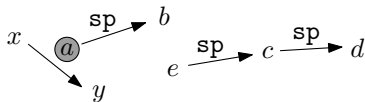
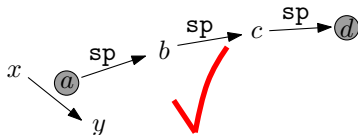
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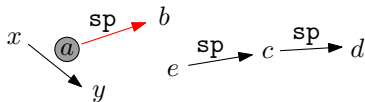
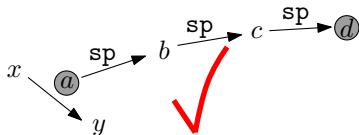
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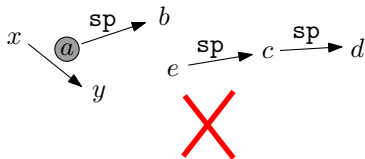
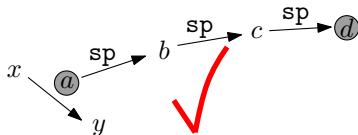
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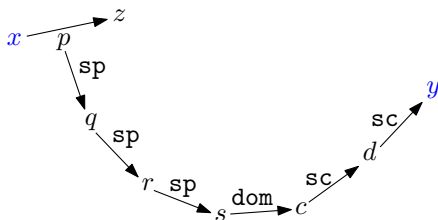
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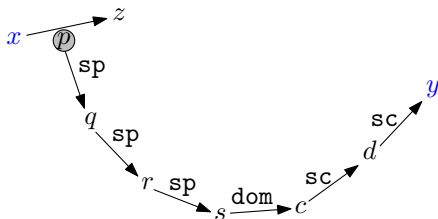
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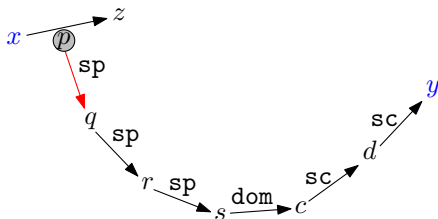
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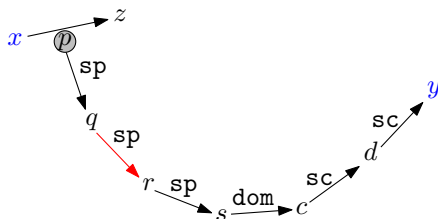
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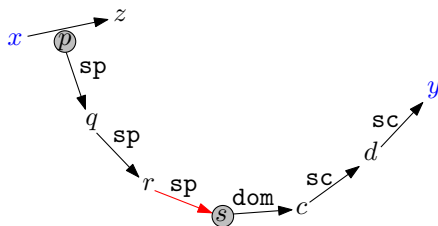
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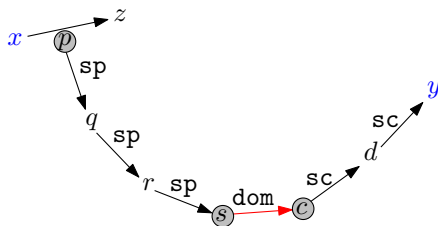
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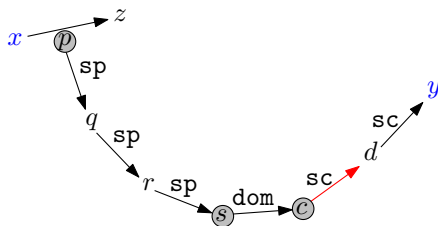
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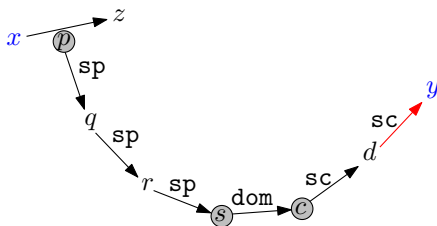
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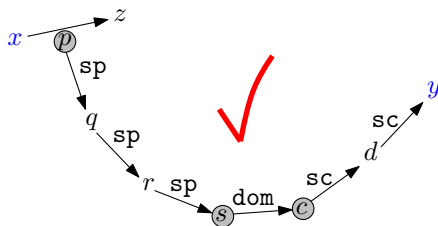
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Theorem

The goal oriented algorithm takes $O(|G| \log |G|)$ time in testing the entailment $G \models t$.

- ▶ Correctness follows by the dependencies diagram.
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The $O(n \log n)$ bound is tight.

Theorem

Testing $G \models t$ takes time $\Omega(|G| \log |G|)$.

Idea: Coding the **set-disjointness** problem, which is $\Omega(n \log n)$

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- ▶ Next: Navigational language based on testing algorithm